

Repa

Regular, Shape-polymorphic, **P**arallel **A**rrays

Gabriele Keller¹ / Manuel Chakravarty¹

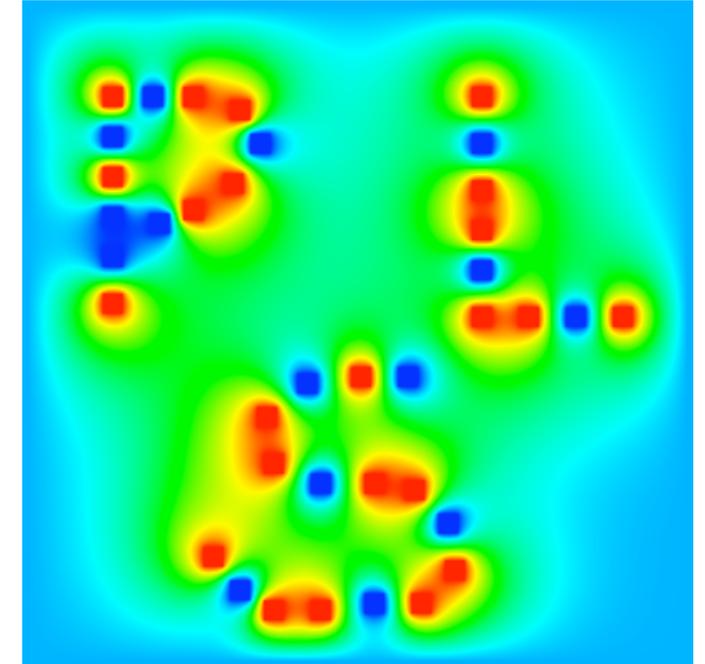
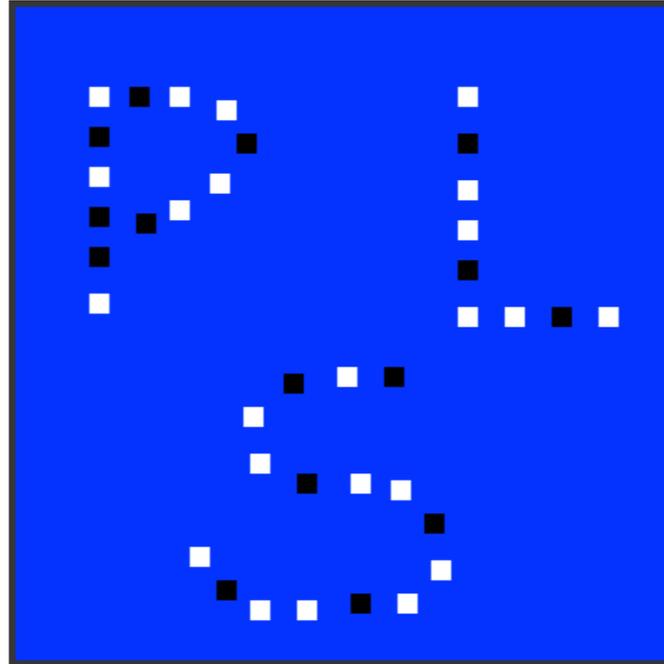
Roman Leshchinskiy¹ / Simon Peyton Jones² / Ben Lippmeier¹

¹University of New South Wales

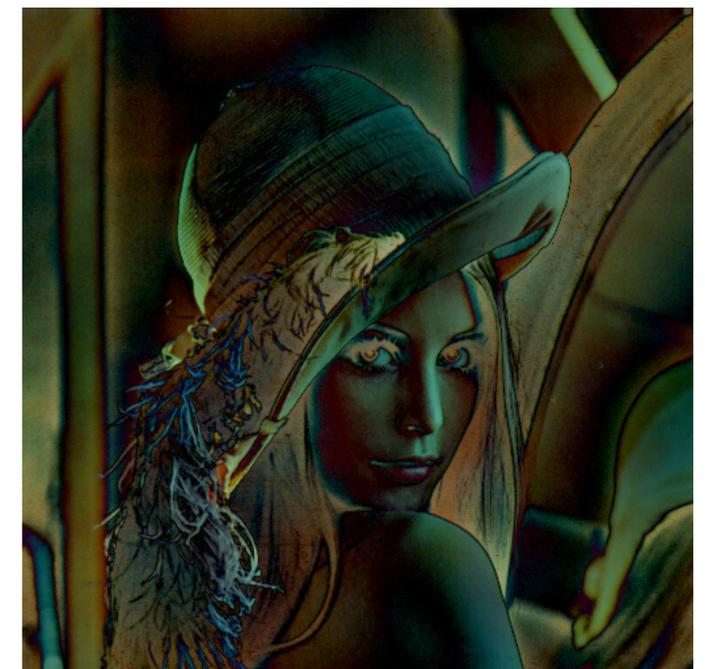
²Microsoft Research Ltd, Cambridge England

Example Applications

Solving the
Laplace Equation

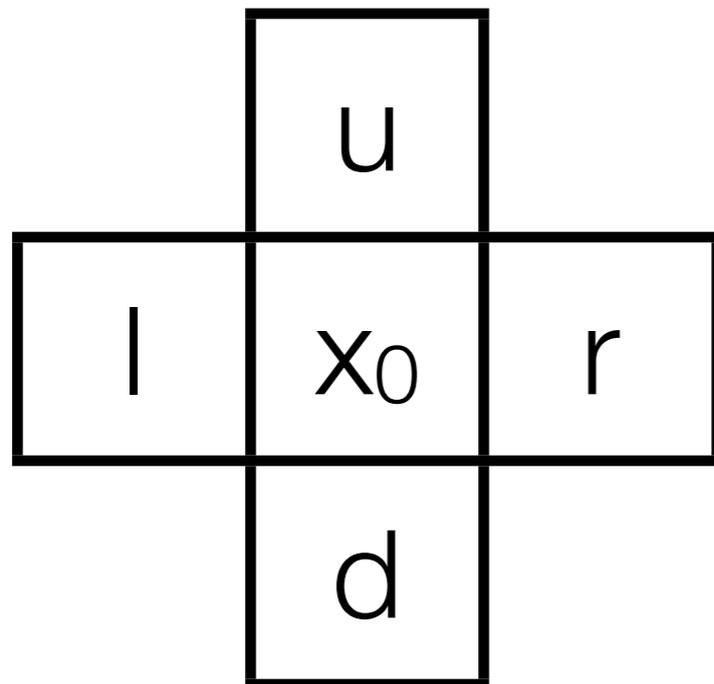
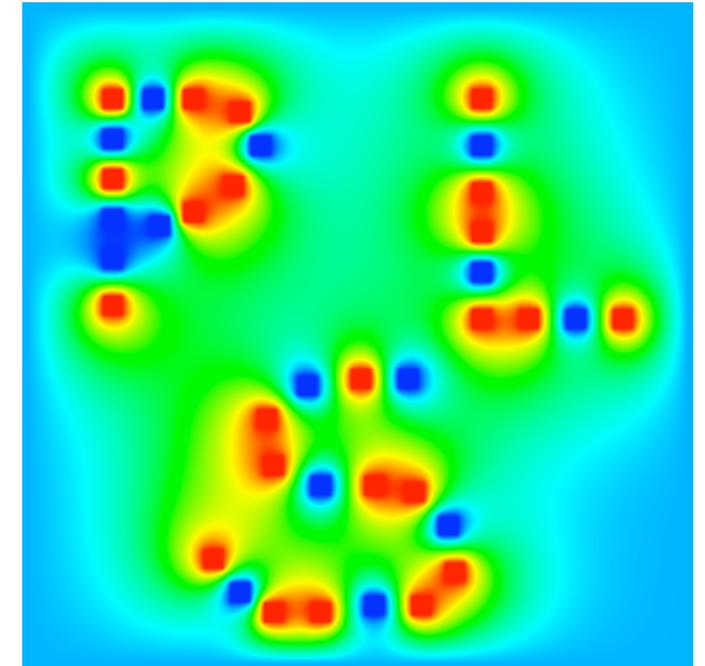
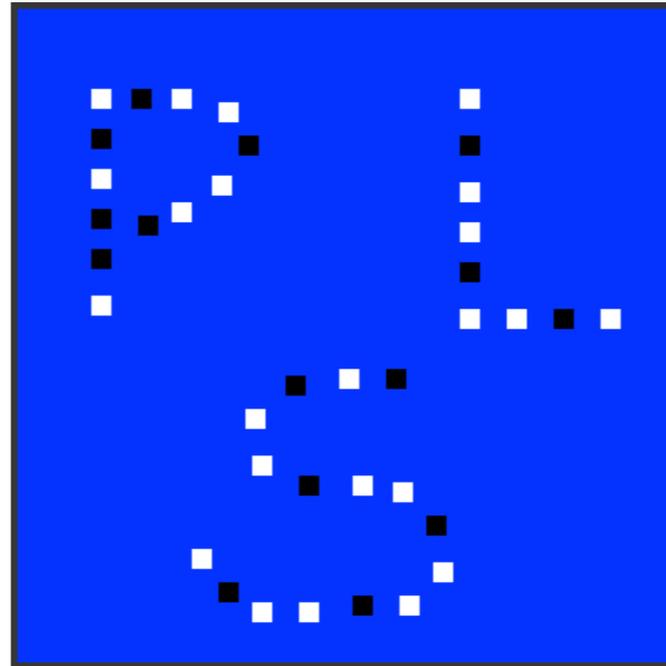


Fast Fourier
Transform
(highpass filter)



Example Applications

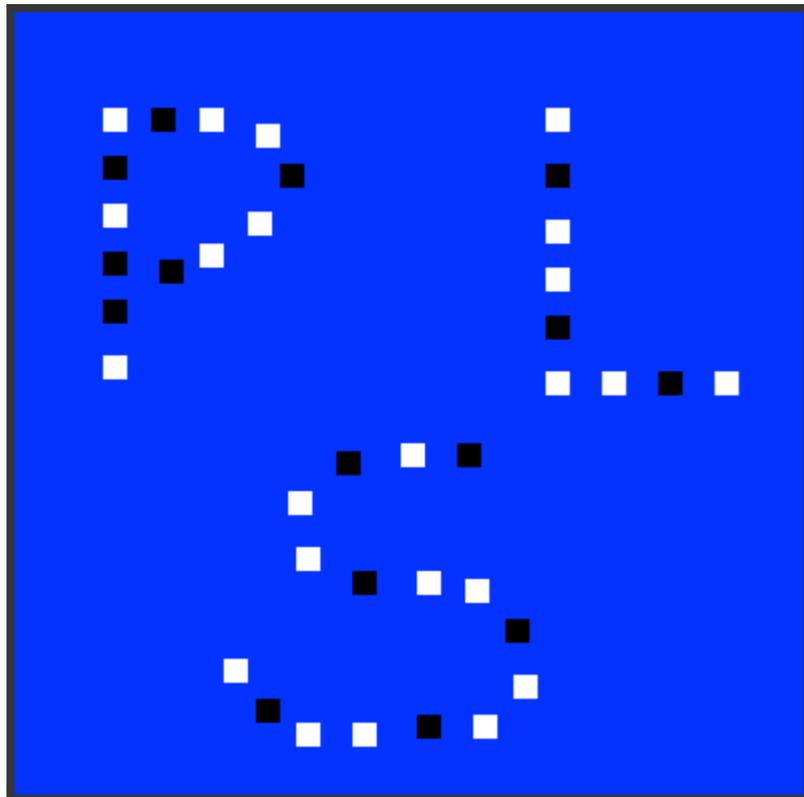
Solving the
Laplace Equation



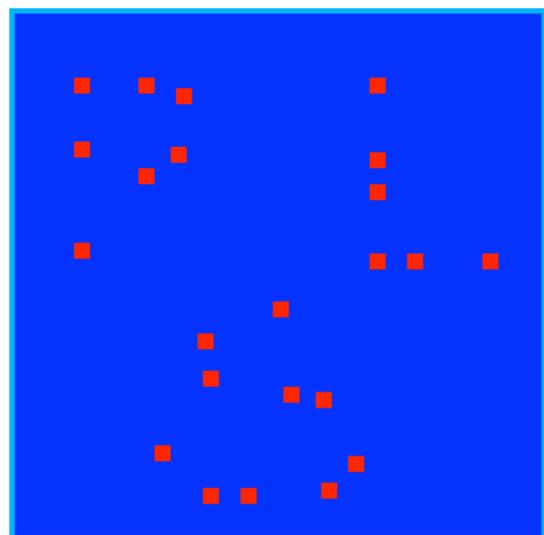
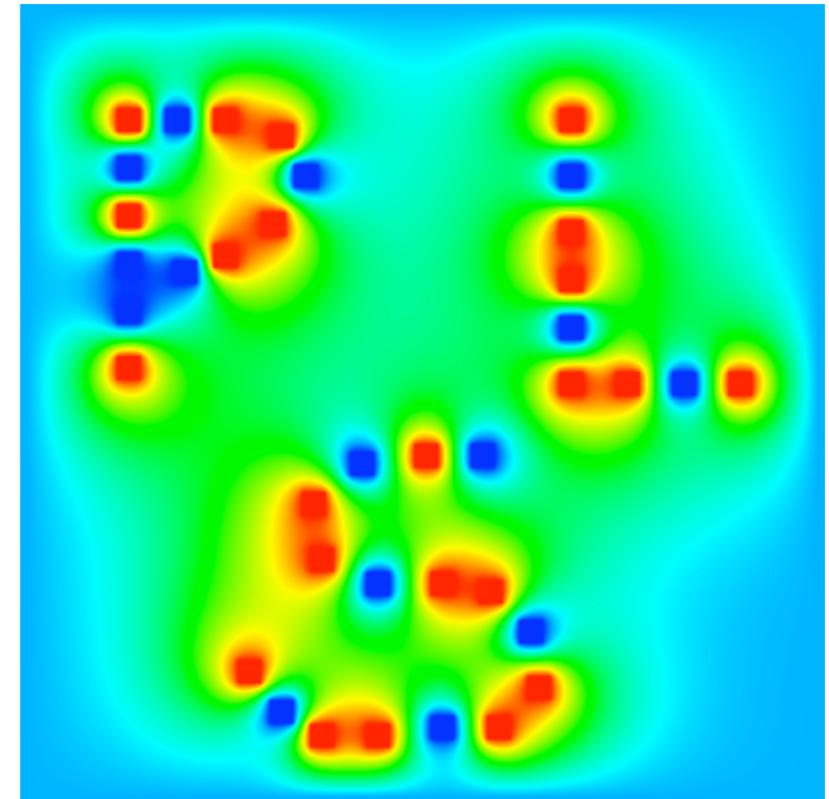
$$x_0' = (l + r + u + d) / 4$$

Laplace Equation

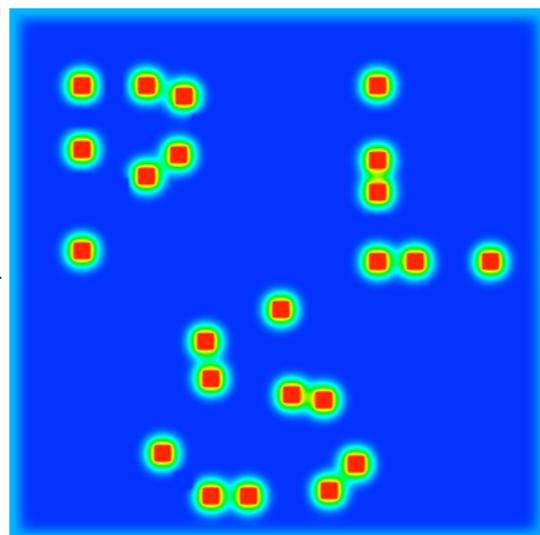
boundary conditions



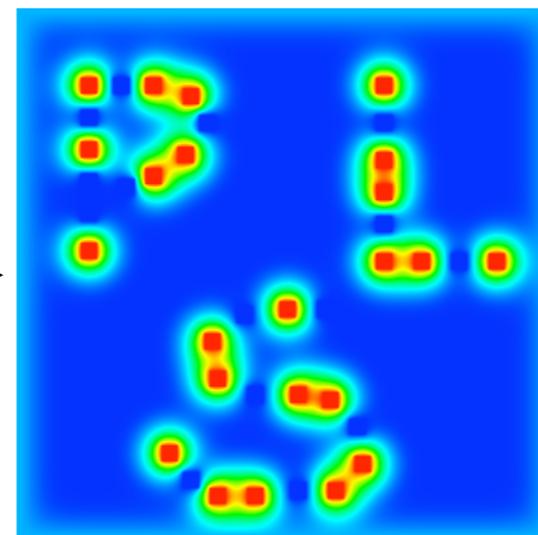
5000 steps



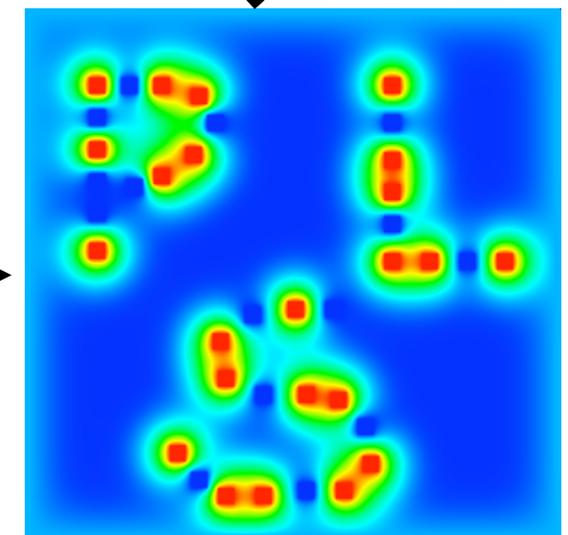
0 steps



100 steps



500 steps



1000 steps

2D Fast Fourier Transform (FFT)

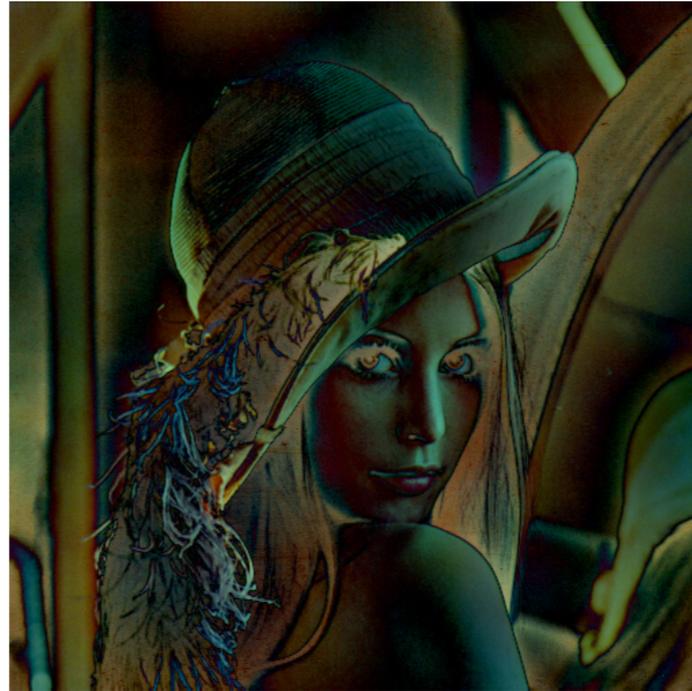
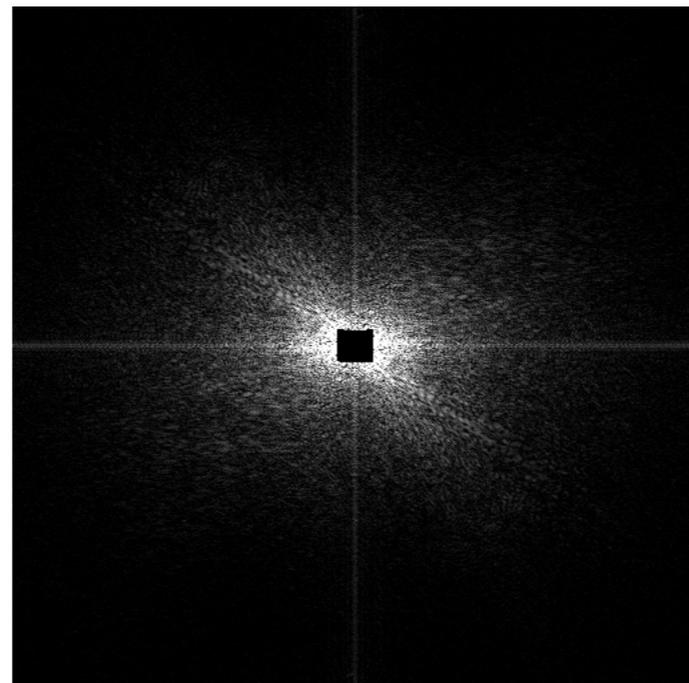
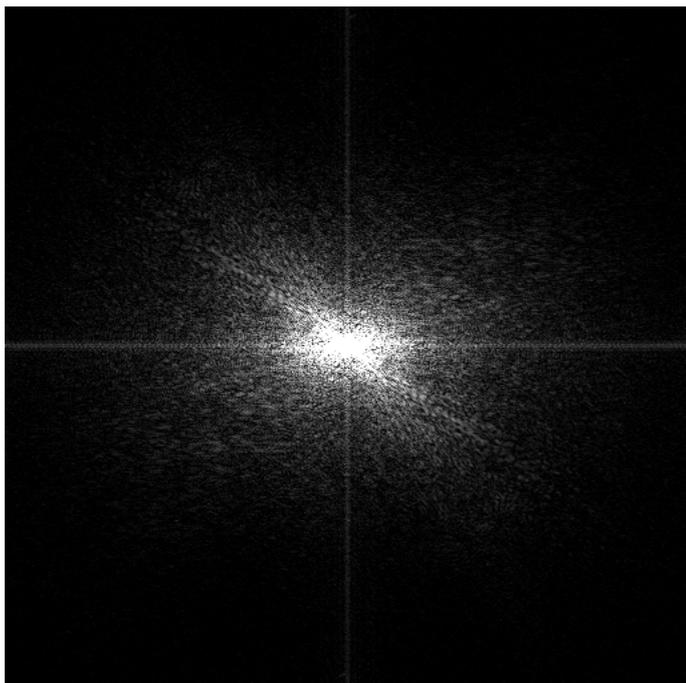


image space

frequency space



Regular, Shape-polymorphic, Parallel Arrays

- **Regular Arrays**

Arrays are dense and rectangular.

Most elements are non-zero.

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Functions work over arrays of arbitrary rank (dimension).

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Most elements are non-zero.

- **Shape Polymorphic**

Functions work over arrays of arbitrary rank (dimension).

- **Flat Data Parallelism**

Individual computations don't need to communicate.

Parallel computations don't spark further parallel computations

Matrix Transposition

```
transpose2D
```

```
:: Elt e => Array DIM2 e -> Array DIM2 e
```

```
transpose2D arr
```

```
= backpermute newExtent swap arr
```

```
where swap (Z :.i :.j) = Z :.j :.i
```

```
newExtent = swap (extent arr)
```

10	20	30
44	55	66

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- An Index Space Transform
- The ordering of the elements changes, but the values do not.

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- An Index Space Transform
- The ordering of the elements changes, but the values do not.
- We usually want to push such transforms into the consumer.

10	20	30
44	55	66

10	44
20	55
30	66

Matrix Multiplication

$$(A \cdot B)_{i,j} = \sum_{k=1}^n A_{i,k} \cdot B_{k,j}$$

a₁₁	a₁₂	a₁₃
a₂₁	a₂₂	a₂₃
a₃₁	a₃₂	a₃₃
a₄₁	a₄₂	a₄₃

•

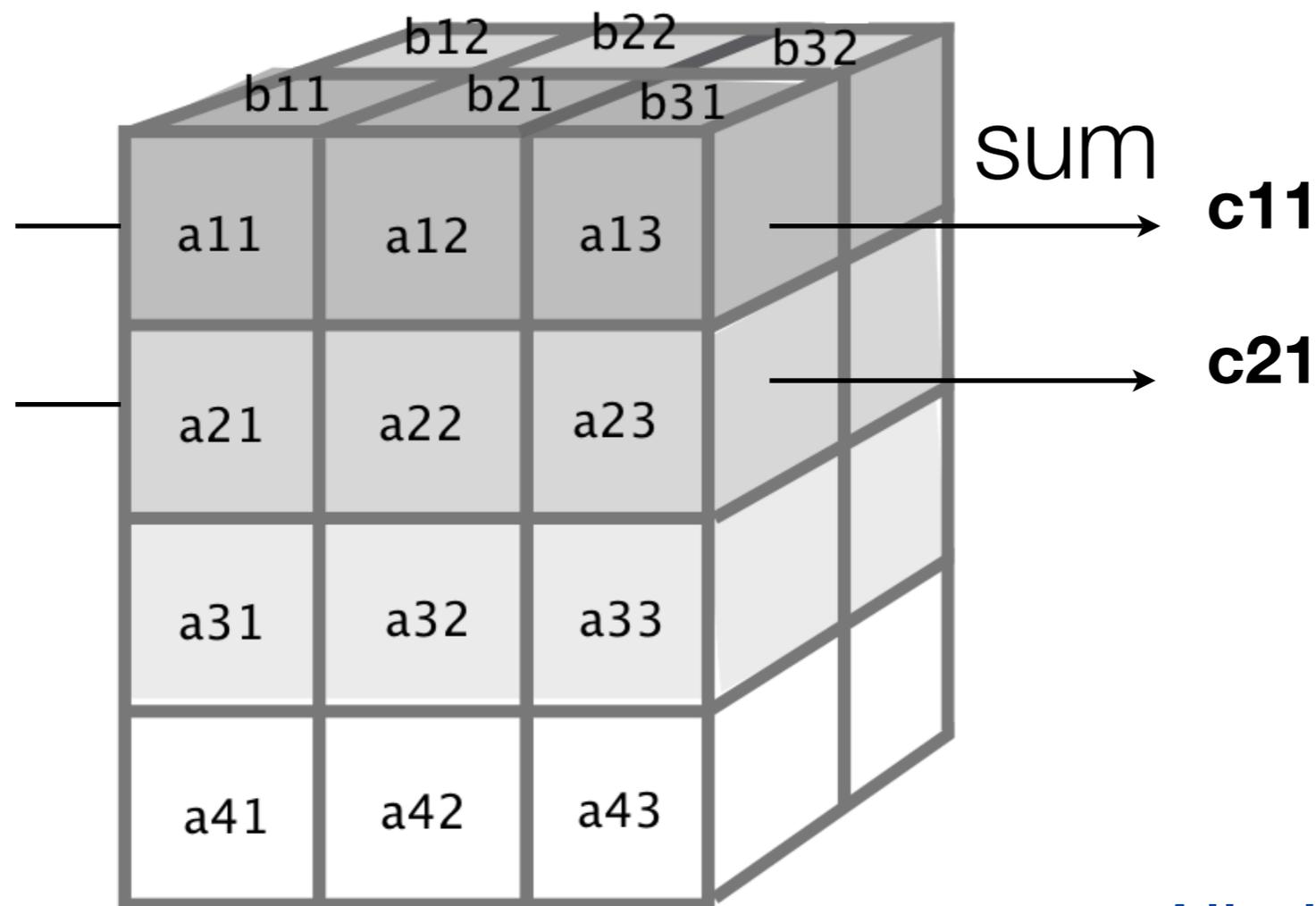
b₁₁	b₁₂
b₂₁	b₂₂
b₃₁	b₃₂

=

c₁₁	c₁₂
c₂₁	c₂₂
c₃₁	c₃₂
c₄₁	c₄₂

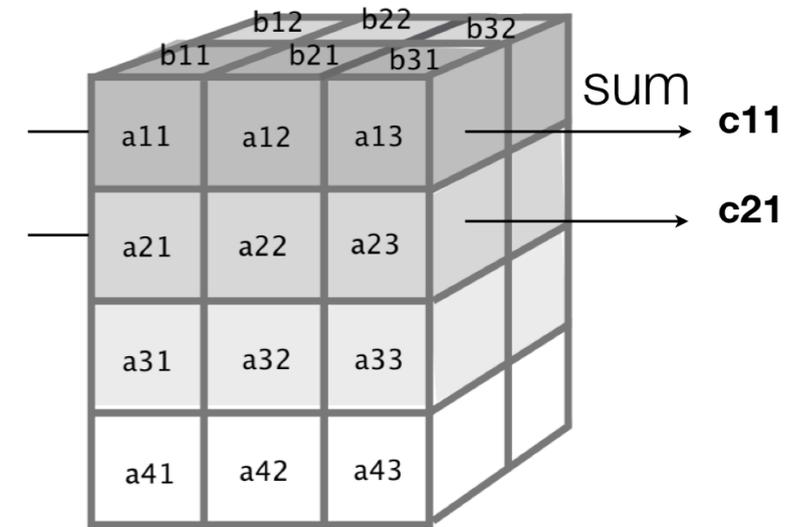
Matrix Multiplication

$$(A \cdot B)_{i,j} = \sum_{k=1}^n A_{i,k} \cdot B_{k,j}$$



- All elements of the result can be computed in parallel!

Matrix Multiplication



```
mmMult
```

```
  :: (Num e, Elt e)
```

```
 => Array DIM2 e
```

```
 -> Array DIM2 e -> Array DIM2 e
```

```
mmMult arr brr
```

```
 = sum (zipWith (*) arrRepl brrRepl)
```

```
where
```

```
  trr = transpose2D brr
```

```
  arrR = replicate (Z :: All    :: colsB :: All) arr
```

```
  brrR = replicate (Z :: rowsA :: All    :: All) trr
```

```
  (Z :: colsA :: rowsA) = extent arr
```

```
  (Z :: colsB :: rowsB) = extent brr
```

Fusion

- It's nice to program with bulk operations
.. but we usually want them to be fused.
- We imagine replicating the source arrays being replicated when writing the program, but we don't want this at runtime.
- Fusion eliminates the intermediate arrays and the corresponding memory traffic.

Manifest and Delayed Arrays

```
data Array sh e
  = Manifest sh (UArr e)
  | Delayed   sh (sh -> e)
```

- **Manifest** wraps a bona-fide unboxed array. Bulk-strict semantics. Forcing one element forces them all.
- **Delayed** wraps an element producing function, perhaps an index transformation that references some other array.
- Delayed functions are inlined and fused by the existing GHC optimiser (and lots of rewrite rules).

Sharing and force

```
let arr = ...  
    brr = map f arr  
in mmMult brr brr
```

Sharing and force

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data Array sh e  
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```

Sharing and force

```
force :: Array sh e  
      -> Array sh e
```

```
data Array sh e  
      = Manifest sh (UArr e)  
      | Delayed   sh (sh -> e)
```

```
let arr = ...  
     brr = force (map f arr)  
in mmMult brr brr
```

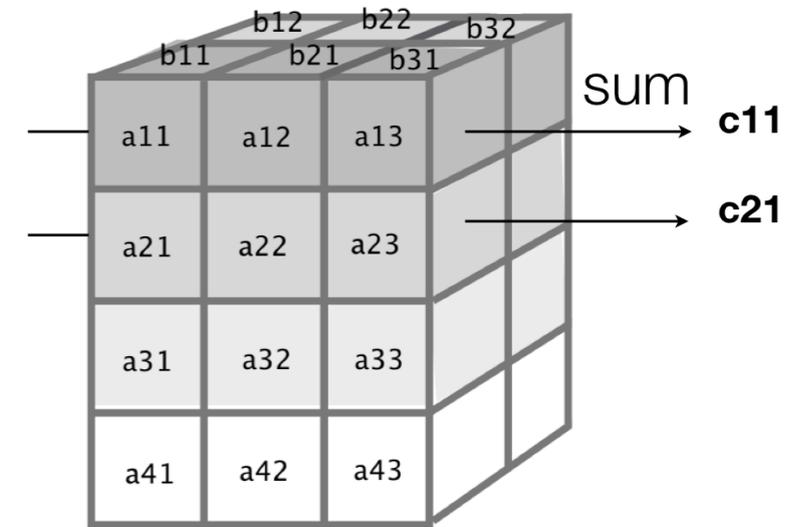
- For **Manifest** arrays, force is the identity.
- For **Delayed** arrays, it evaluates all the elements in parallel, producing a manifest array.
- The programmer must add force manually.

Using the force ...

$$\begin{array}{|c|c|c|} \hline \mathbf{a}_{11} & \mathbf{a}_{12} & \mathbf{a}_{13} \\ \hline \mathbf{a}_{21} & \mathbf{a}_{22} & \mathbf{a}_{23} \\ \hline \mathbf{a}_{31} & \mathbf{a}_{32} & \mathbf{a}_{33} \\ \hline \mathbf{a}_{41} & \mathbf{a}_{42} & \mathbf{a}_{43} \\ \hline \end{array} \cdot \begin{array}{|c|c|} \hline \mathbf{b}_{11} & \mathbf{b}_{12} \\ \hline \mathbf{b}_{21} & \mathbf{b}_{22} \\ \hline \mathbf{b}_{31} & \mathbf{b}_{32} \\ \hline \end{array} = \begin{array}{|c|c|} \hline \mathbf{c}_{11} & \mathbf{c}_{12} \\ \hline \mathbf{c}_{21} & \mathbf{c}_{22} \\ \hline \mathbf{c}_{31} & \mathbf{c}_{32} \\ \hline \mathbf{c}_{41} & \mathbf{c}_{42} \\ \hline \end{array}$$

- We get better cache performance when accessing the **b** elements left-to-right rather than top-to-bottom

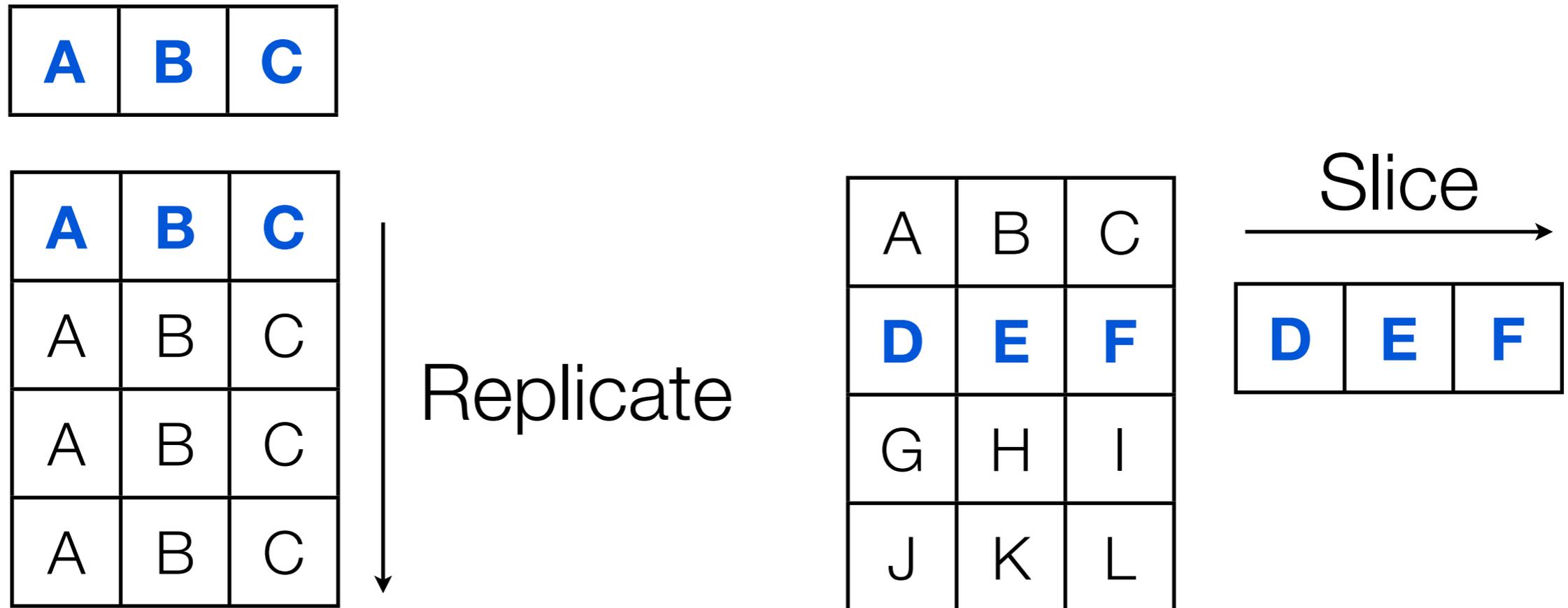
Using the force ...



```
mmMult
  :: (Num e, Elt e)
  => Array DIM2 e
  -> Array DIM2 e -> Array DIM2 e
```

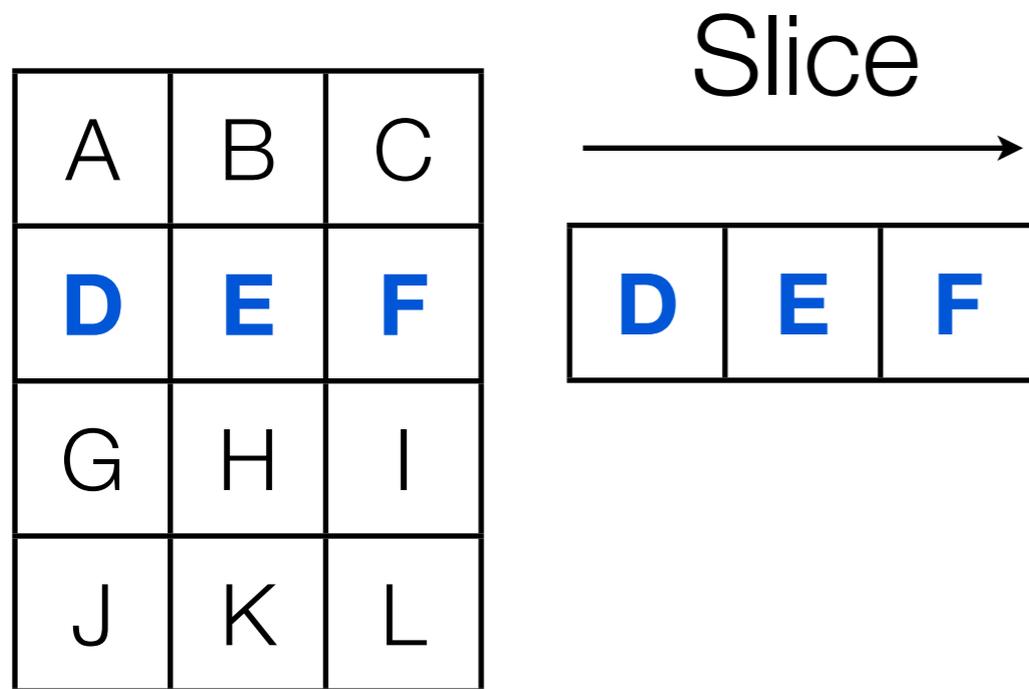
```
mmMult arr brr
= sum (zipWith (*) arrRepl brrRepl)
where
  trr = force (transpose2D brr)
  arrR = replicate (Z :: All :: colsB :: All) arr
  brrR = replicate (Z :: rowsA :: All :: All) trr
  (Z :: colsA :: rowsA) = extent arr
  (Z :: colsB :: rowsB) = extent brr
```

Replicate and Slice are duals.



- Replicate and Slice are index transforms.
- The values of the array elements do not change.

Type hackery



Type hackery

A	B	C
D	E	F
G	H	I
J	K	L

Slice

→

D	E	F
----------	----------	----------

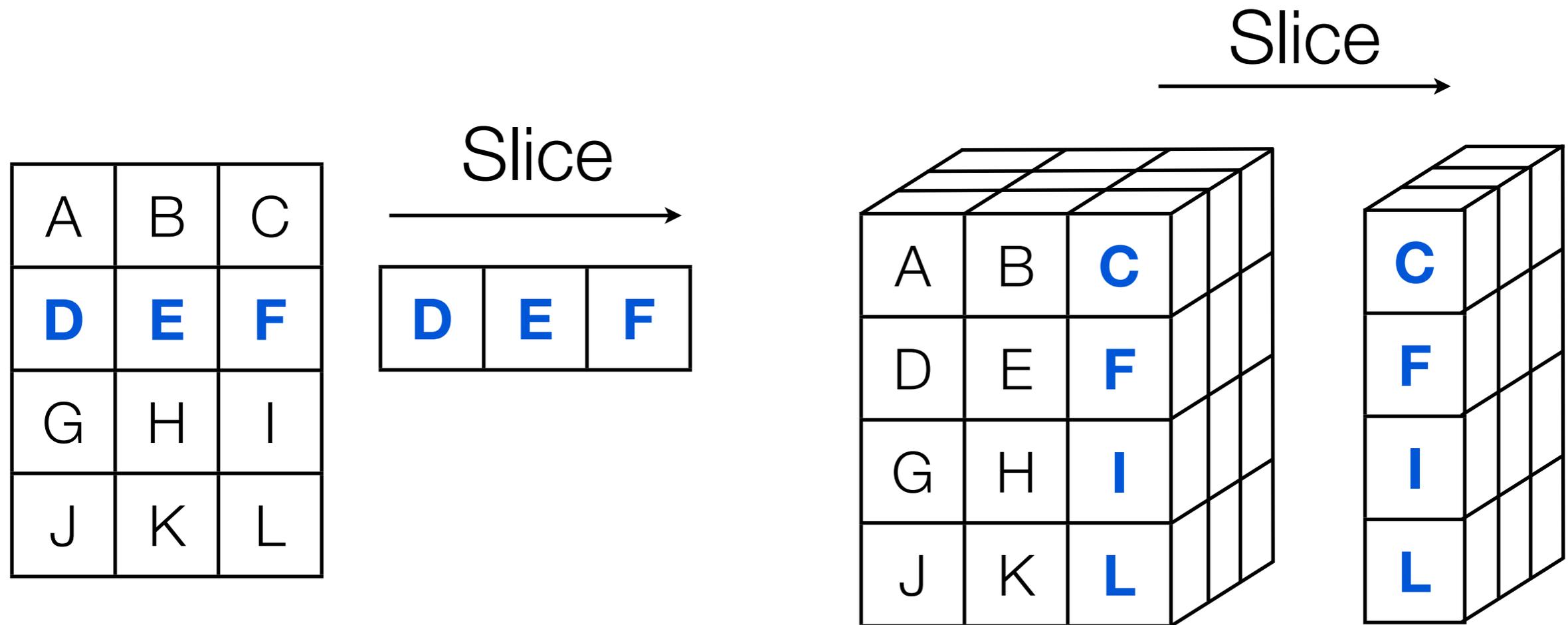
Slice

→

A	B	C
D	E	F
G	H	I
J	K	L

C
F
I
L

Type hackery



```
slice :: ( Slice s1, Elt e
         , Shape (FullShape  s1))
         , Shape (SliceShape s1))
=> Array (FullShape s1)  e
-> s1 -> Array (SliceShape s1) e
```

Other operations

map :: (Shape sh, Elt a, Elt b)
 => (a -> b) -> Array sh a -> Array sh b

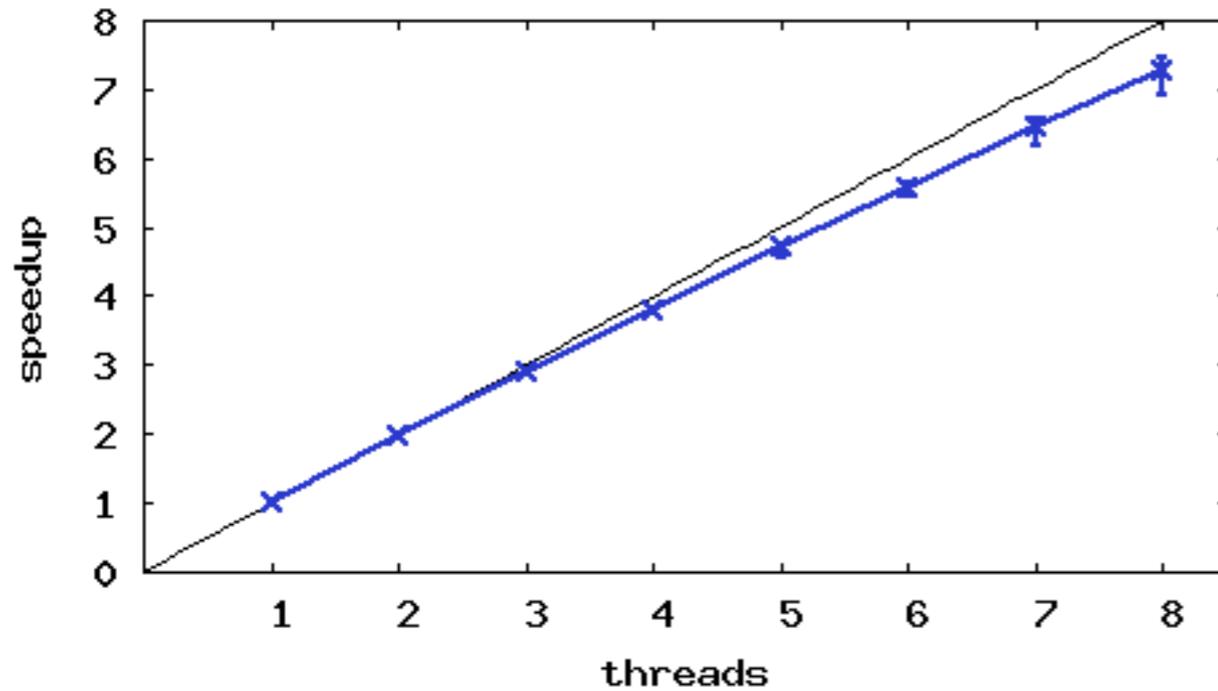
zip :: (Shape sh, Elt a, Elt b)
 => Array sh a -> Array sh b
 -> Array sh (a, b)

foldl :: (Shape sh, Elt a, Elt b)
 => (a -> b -> a) -> a
 -> Array (sh :: Int) e -> Array sh a

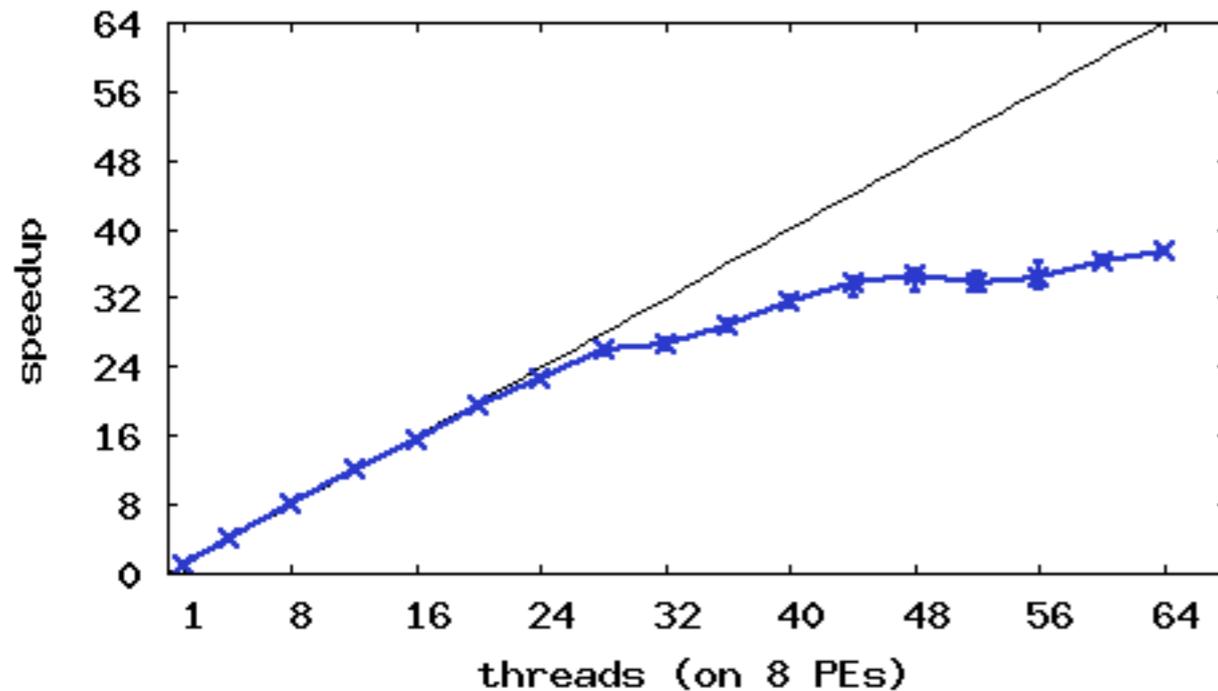
reshape :: (Shape sh, Shape sh', Elt e)
 => sh -> Array sh' e -> Array sh e

Matrix Multiplication 1024x1024

on a 2x Quad-core 3Ghz Xenon



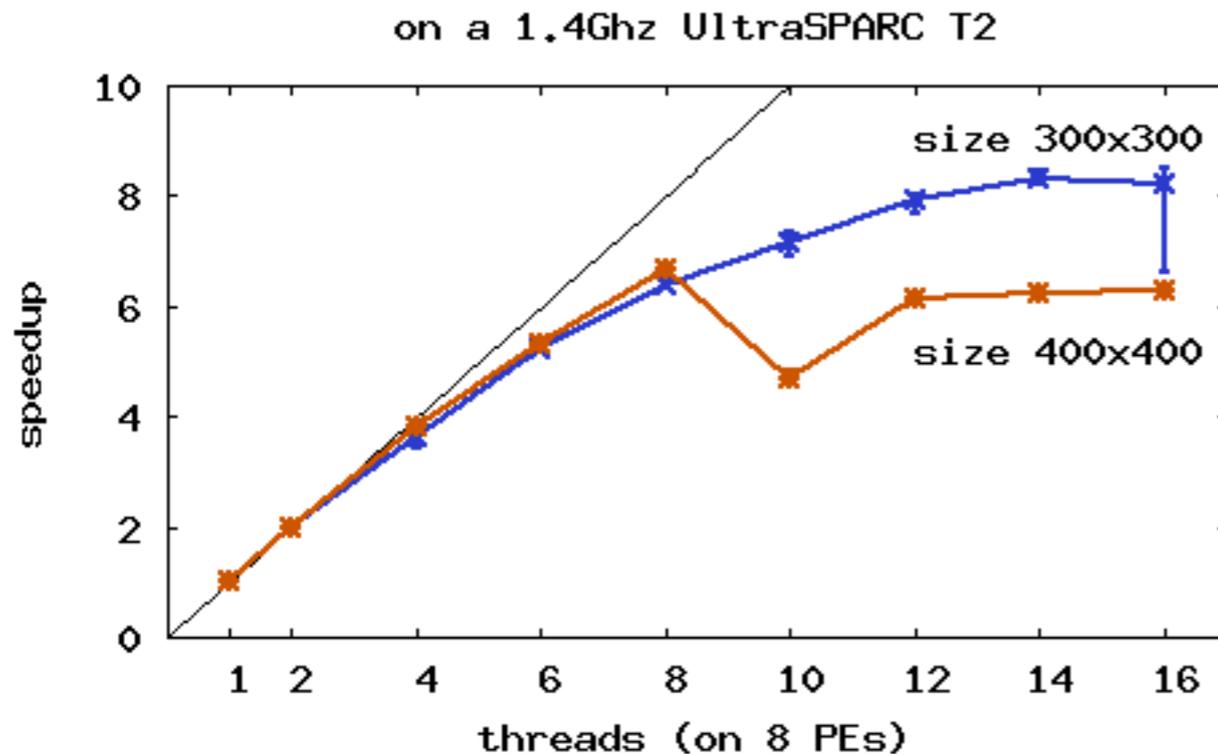
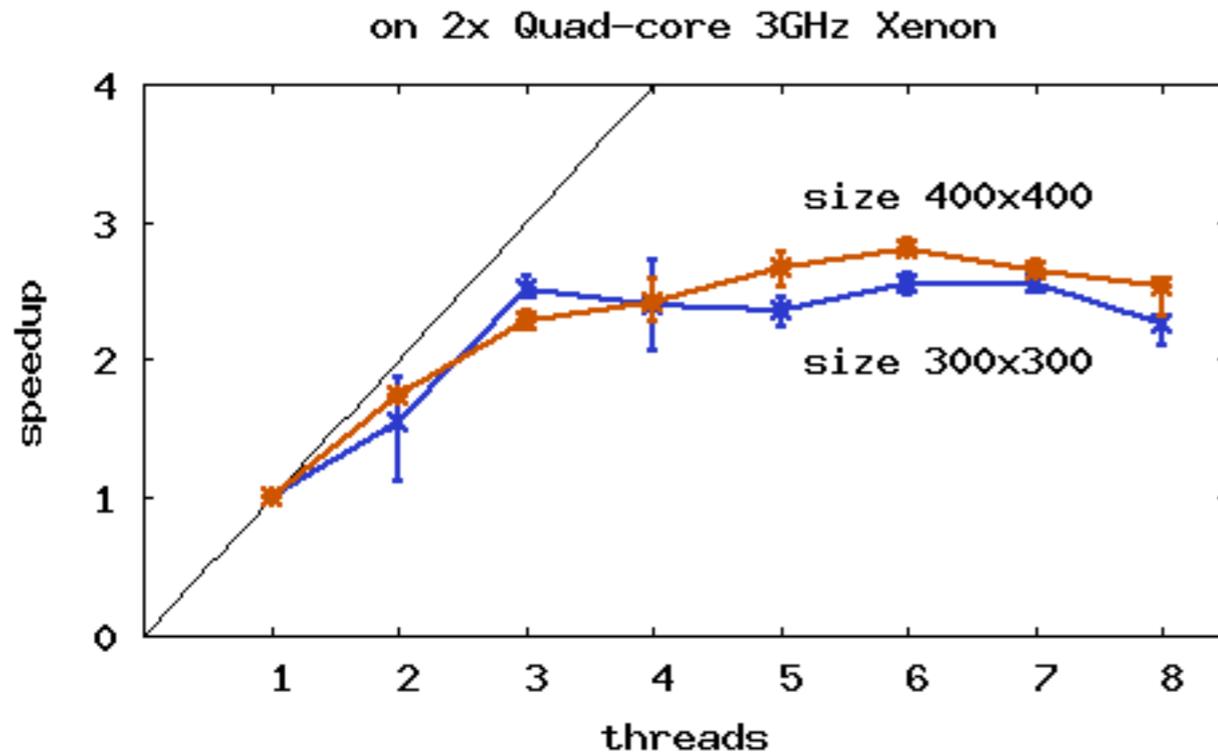
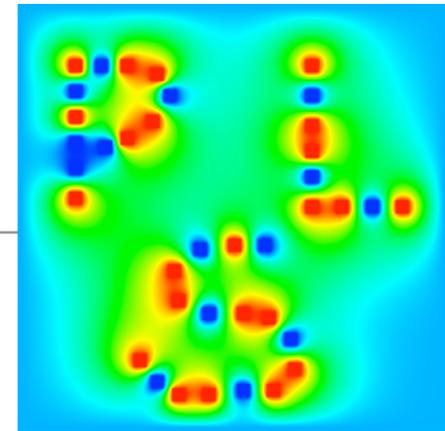
on a 1.4Ghz UltraSPARC T2



	GCC	single thread	fastest parallel
Xenon	3.8s	4.6s	0.64s
T2	52s	92s	2.4s

- C reference version uses double nested loops.
- Exposing sufficient parallelism on the T2 is a must.

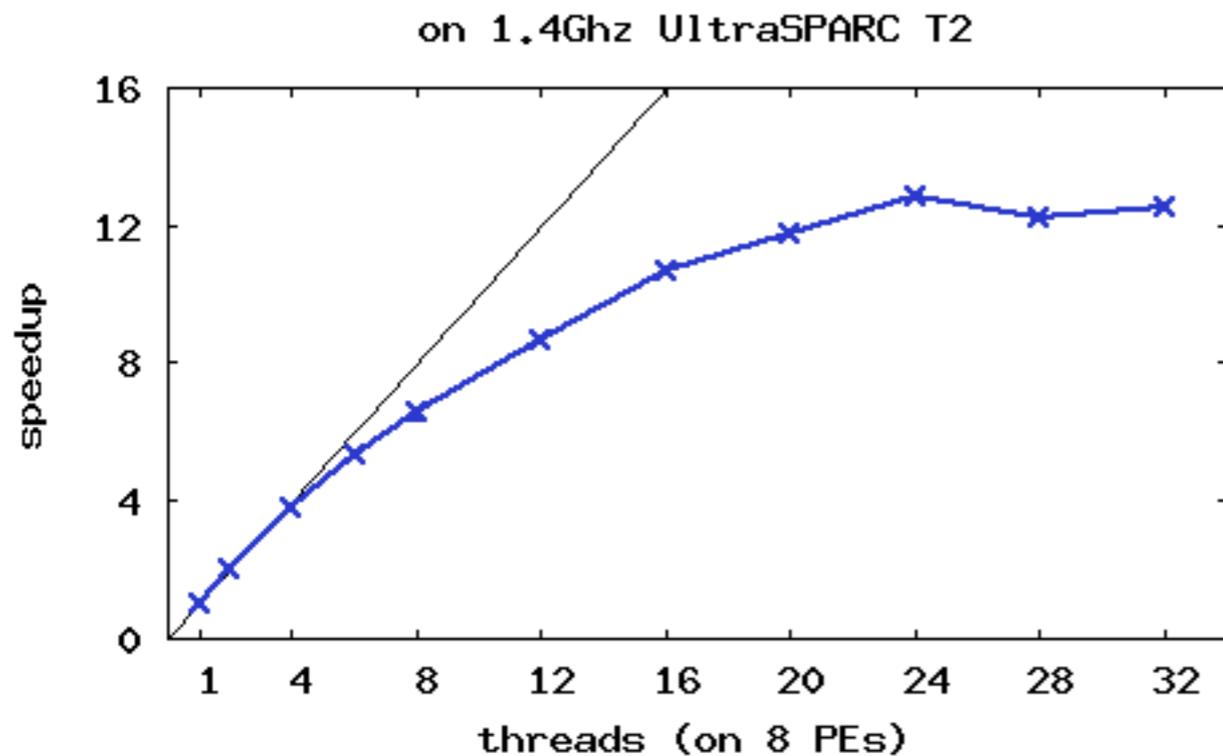
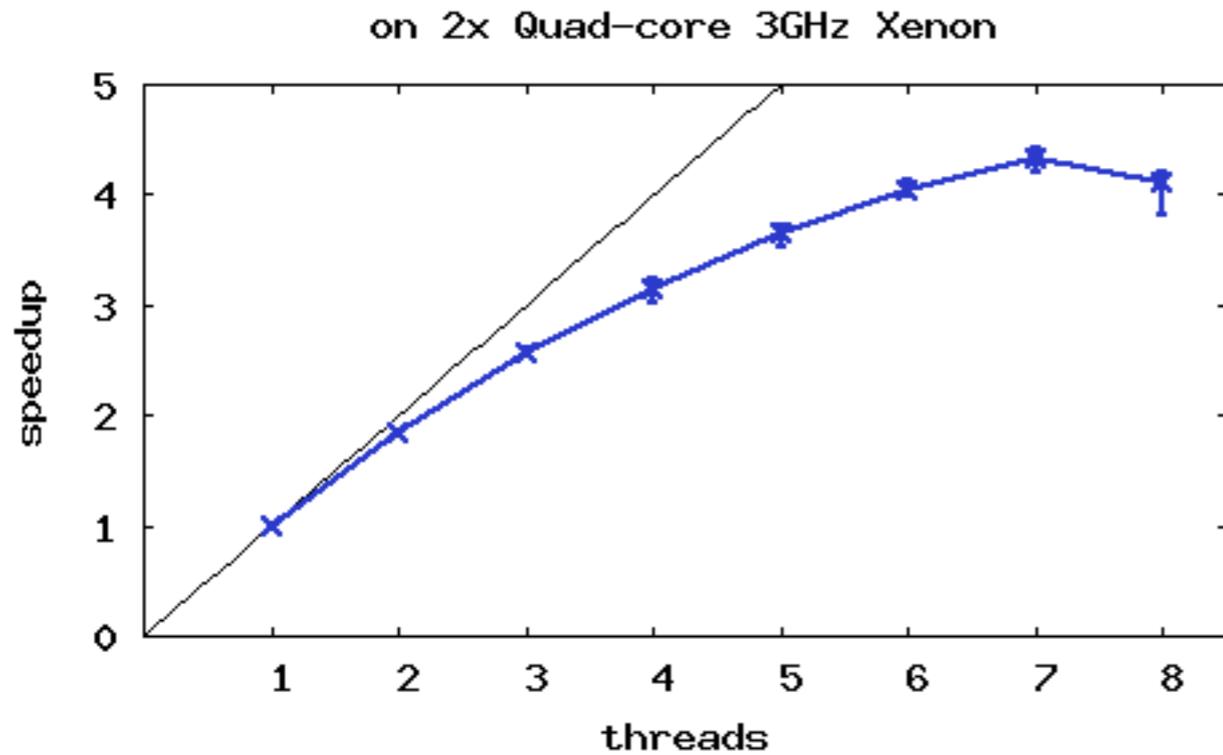
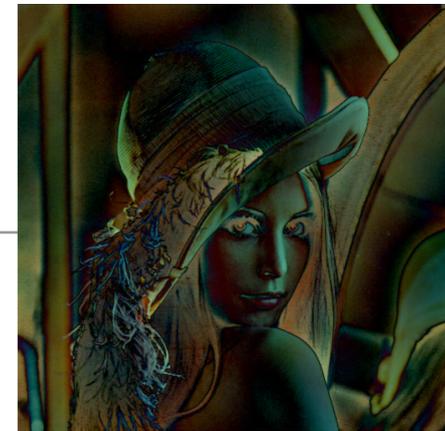
Laplace Equation



	GCC	single thread	fastest parallel
Xenon	0.70	1.7s	0.68s
T2	6.5s	32s	3.8s

- GHC native code generator does no instruction reordering on SPARC. No LLVM 'port.
- Single threaded on T2 is slow

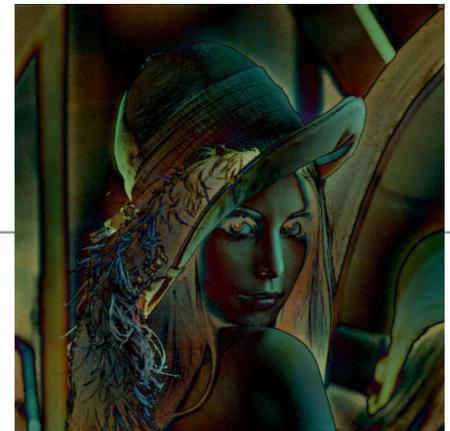
2D Fast Fourier Transform (FFT)



	GCC	single thread	fastest parallel
Xenon	0.24	8.8s	2.0s
T2	2.4s	98s	7.7s

- C version is FFTW which uses in-place deep magic.
- Parallelism is no substitute for a better algorithm.

2D Fast Fourier Transform (FFT)



```
fft1D :: Array (sh::Int) Double
      -> Array (sh::Int) Double
      -> Array (sh::Int) Double
```

```
fft1D rofu
```

```
| n > 2 = (left +^ right) :+: (left -^ right)
```

```
| n == 2 = traverse v id swivel
```

where

```
(_ :: n) = extent v
```

```
swivel f (ix:.0) = f (ix:.0) + f (ix:.1)
```

```
swivel f (ix:.1) = f (ix:.0) - f (ix:.1)
```

```
rofu' = evenHalf rofu
```

```
left = force . (fft1D rofu' . evenHalf $ v
```

```
right = force . (*^ rofu) . fft1D rofu' . oddHalf $ v
```

Future work

- The examples we've presented are easy to write, but are cache naive.

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Future work

- The examples we've presented are easy to write, but are cache naive.
- Using, block based matrix multiplication imposes a restriction on the order of evaluation...
... and makes it less obvious how to parallelise.
- Repeated index computations can be expensive.
GHC does not perform strength reduction on its loops.

Get it

<http://trac.haskell.org/repa>

- We depend on the the current version of the GHC head for decent performance.
- There will be a new release in a few weeks with GHC 7.0
- Send me your programs and I'll add them to our performance regression testsuite!

Questions?

Spare Slides

Shapes and Indices

```
data Z = Z
data tail :: head = tail :: head

type DIM0 = Z
type DIM1 = DIM0 :: Int
type DIM2 = DIM1 :: Int
...

class Shape sh where
  rank    :: sh -> Int
  size    :: sh -> Int
  toIndex :: sh -> sh -> Int
  fromIndex :: sh -> Int -> sh

instance Shape Z where ...
instance Shape sh => Shape (sh :: Int) where ...
```

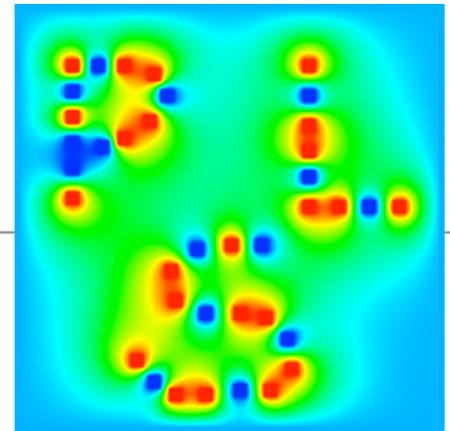
Generic Traversal and the 3rd order function

traverse

```
  :: (Shape sh, Shape sh', Elt e)
=> Array sh e
-> (sh -> sh')           -- shape transform
-> ((sh -> e) -> sh' -> e') -- elem transform
-> Array sh' e'
```

- A sane use for a third order function!
- Traverse takes a function to calculate the elements of the array.
- That function is passed a function to get elements of the source array.

Laplace Equation



```
stencil :: Array DIM2 Double  
        -> Array DIM2 Double
```

```
stencil arr  
= traverse arr id update
```

where

```
_ :: height :: width = extent arr
```

```
update get d@(sh :: i :: j)  
= if isBoundary i j  
  then get d  
  else (get (sh :: (i-1)) :: j)  
       + get (sh :: i :: (j-1))  
       + get (sh :: (i+1)) :: j)  
       + get (sh :: i :: (j+1))) / 4
```